

# BRANCH-PRICE-AND-CUT FOR CAUSAL DISCOVERY James Cussens, Dept of Computer Science, University of Bristol

## **1. Casual discovery and optimisation**

- There are pros and cons to doing score-based causal discovery (whatever the class of candidate causal models) using a general-purpose constrained optimisation solver.
- Pro: The (highly optimised) software is already there.
- Pro: Can easily add constraints to rule out causal models inconsistent with domain knowledge
- Pro: Exact (sometimes) and anytime (always) learning is possible.
- Con: Need to encode the causal discovery problem in a way the chosen solver understands
- Con: More complex than using a causal-discovery-specific algorithm.

## 2. Encoding DAGs as vectors

This DAG 
$$i \xrightarrow{j}$$
 is this vector in  $\mathbb{R}^{12}$ :

$x_{i \leftarrow \{\}}$	$x_{i \leftarrow \{j\}}$	$x_{i \leftarrow \{k\}}$	$x_{i \leftarrow \{j,k\}}$	$x_{j \leftarrow \{\}}$	$x_{j \leftarrow \{i\}}$	$x_{j \leftarrow \{k\}}$	$x_{j \leftarrow \{i,k\}}$	$x_{k \leftarrow \{\}}$	$x_{k \leftarrow \{i\}}$	$x_{k \leftarrow \{j\}}$	$x_{k \leftarrow \{i,j\}}$
0	1	0	0	1	0	0	0	0	0	0	1

• Why this encoding? Because many objective functions ('scores') for DAGs are sums of *local scores* which are determined by the choice of *parents* for each vertex.

### **3. Integer linear programming model for DAG learning**



SUBJECT TO:

(is a directed graph) 
$$\sum_{J \subseteq P \setminus \{i\}} x_{i \leftarrow J} = 1 \qquad i \in P \quad (1)$$
  
(is acyclic) 
$$\sum_{i \in C} \sum_{\substack{J \subseteq P \setminus \{i\}\\ J \cap C \neq \emptyset}} x_{i \leftarrow J} \leq |C| - 1 \qquad C \subseteq P, |C| \geq 2 \quad (2)$$

 $x_{i\leftarrow J} \in \{0,1\}, \qquad i \in P, J \subseteq P \setminus \{i\}$ 

- P are the random variables=vertices of the DAG.
- $x_{i \leftarrow J}$  indicates that J is the parent set for child i.
- $c_{i \leftarrow J}$  is the local cost (-1 × local score) when J is the parent set for child *i*.

# 4. Cutting and pricing to solve large linear programs

- The initial step for an ILP solver is to solve the *linear relaxation* of the ILP problem. In our case, this linear relaxation is a linear program (LP) where  $x_{i\leftarrow J} \in \{0, 1\}$  is replaced  $x_{i\leftarrow J} \in [0, 1]$ .
- There are  $2^p p 1$  (p = |P|) cluster constraints ruling out cycles so we add them as cutting planes. We initially solve an LP (LP<sub>0</sub>) with no cluster constraints to get a solution  $x_0^*$  and then add only those cluster constraints that  $x_0^*$  violates to get a new LP (LP<sub>1</sub>), and then resolve (to get solution  $x_1^*$ ). We repeat until we get an LP (LP<sub>m</sub>) whose solution  $x_m^*$  satisfies all cluster constraints (even though only a fraction of them are included in the problem).
- But each LP in the sequence  $LP_0, LP_1, \ldots, LP_m$  has  $p2^{p-1}$  variables!
- For each  $LP_{\iota}$  we only need to include those  $x_{i \leftarrow J}$  variables that have a non-zero value in the solution  $x_{\iota}^*$ . (An ILP variable not included in the problem is implicitly set to 0.)
- So, once we have solved  $LP_{\iota}$  with the ILP variables currently in the problem, we look for additional ILP variables which, if included, would lead to a better (lower cost) solution of  $LP_{\iota}$ . If we can't find any then  $LP_{\iota}$  has been solved to optimality even though typically a small fraction of ILP variables have been included. This is called *pricing*.

6. Pricing for  $\ell_0$  penalised Gaussian DAGs

- Associated with each LP solution  $x_{\iota}^*$  there are *dual values* for each constraint in LP<sub> $\iota$ </sub>.
- Let  $\lambda_i^*$  be the dual value for the constraint (1) for child *i* and let  $\lambda_C^*$  be the dual value for the constraint (2) for cluster *C*.
- A variable  $x_{i \leftarrow J}$  is worth adding to the problem if its reduced cost  $c_{i \leftarrow J} \lambda_i^* + \sum_{\substack{C \in \mathcal{C}, i \in C \\ C \cap J \neq \emptyset}} \lambda_C^*$  is negative.
- We search for new variables whose reduced cost is minimal:

$$\label{eq:minimize} \begin{split} \text{MINIMISE}_J \qquad & z = c_{i \leftarrow J} - \lambda_i^* + \sum_{\substack{C \in \mathcal{C}, i \in C \\ C \cap J \neq \emptyset}} \lambda_C^* \\ \text{SUBJECT TO} \qquad & z < 0 \end{split}$$

• For Gaussian DAGs, when the cost is negative log-likelihood with an  $\ell_0$  penalty the pricing problem becomes:

$$\begin{split} \text{MINIMISE}_J \qquad z = n \log \sigma_{i \leftarrow J}^2 + \Lambda^2 |J| + \sum_{\substack{C \in \mathcal{C}, i \in C \\ C \cap J \neq \emptyset}} \lambda_C^* \\ \text{SUBJECT TO} \qquad z < \lambda_i^* \end{split}$$

• n is the size of the data.  $\Lambda^2$  is the  $\ell_0$  penalty.

- $\sigma_{i \leftarrow J}^2$  is the MSE (with MLE parameters) for the linear regression model where variables J predict child i.
- Note this is doubly penalised regression: we have the normal  $\ell_0$  penalty  $\Lambda^2 |J|$ , but also a 'cyclicity' penalty  $\sum_{\substack{C \in \mathcal{C}, i \in C \\ C \cap J \neq \emptyset}} \lambda_C^*$

## 7. Optimisations

5. Pricing

• If a candidate parent set J for child i has a subset J' with lower cost then it is NOT a *potentially optimal parent set (POP)* and we can rule out J as a parent set for i:

 $\forall i \in P, J \subseteq P \setminus \{i\} : \exists J' \subsetneq J : c_{i \leftarrow J} \ge c_{i \leftarrow J'} \to x_{i \leftarrow J} = 0$ 

• Non-POPs should not be priced-in, even if they have negative reduced costs.

## 8. Implementation and Performance

- GOBNILP has been extended to allow pricing to introduce new ILP variables during solving.
- GOBNILP uses SCIP which (unlike e.g. Gurobi) has support for pricing.
- The pricing problem for the Gaussian log-likelihood  $\ell_0$ -penalised score is solved by SCIP
- We can identify set intervals  $\{J : \underline{J} \subseteq J \subseteq \overline{J}\}$  that do not contain POPs and tell the pricing algorithm not to search there.
- Another simpler optimisation is to find all POPs up to some small cardinality before solving begins and tell the pricer to only look for bigger parent sets.
- We can also *delay pricing* so that we only price in new variables once no further cutting planes can be found.
- In one 20 BN variable learning task delayed pricing reduced (exact) solving from 2206 seconds to 351 seconds.
- solving the non-linear optimisation problem (given above).
- If the optimal DAG is sparse then it is often possible to find all POPs before solving starts. This is much faster than pricing them in during solving.
- However, when the optimal DAG is dense, pricing is necessary.
- In one example, produced after the paper was submitted (!), where the optimal DAG had 20 vertices with only one non-adjacency, GOBNILP-with-pricing found it (and proved it optimal) in 140 seconds.
- But standard GOBNILP ran out of memory after 9 minutes.

### 9. Conclusions

More work is required to determine the usefulness and limitations of Branch-Price-and-Cut for Causal Discovery. It would be interesting to apply it to causal discovery beyond DAG learning under causal sufficiency (e.g. to MAG learning).